

PHYSICS 391
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QUANTUM INFORMATION THEORY I

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1 Introduction

A method is more important than a discovery, since the right method will lead to new and even more important discoveries.

Lev Landau (1908-1968)

The objective of this course was to introduce ourselves to quantum information theory, a booming topic in physics with tremendous practical applications. This paper is based on a lecture I gave for the Kapitza Society.

In the first section we will learn about the density operator: a powerful computing tool in quantum mechanics. We will then have a brief overview of Lagrange multipliers and how to use them to maximize entropy and retrieve the canonical ensemble's properties. Finally, we will reveal a surprising relation between statistical and quantum mechanics.

2 The Density Operator

2.1 Quantum Mechanical Ensembles

We generally may not have perfect knowledge of a prepared quantum state. Suppose a third party, Bob, prepares a state for us and only gives a probabilistic description of it, i.e., Bob selects $|j\rangle_x$ with probability $p_X(x)$, where p_X is the probability distribution for the random variable X , and x is in some alphabet \mathcal{X} . We can summarize this information by defining an ensemble E of quantum states

$$E = \{p_X(x); |j\rangle_x\}_{g_x} \quad (1)$$

For example, $E = \{\frac{1}{3}; |j1\rangle; \frac{2}{3}; |j3\rangle\}$. If this ensemble is span by $\{|j0\rangle; |j1\rangle; |j2\rangle; |j3\rangle\}$, then the state is $|j0\rangle$ with probability 0, $|j1\rangle$ with probability $\frac{1}{3}$, $|j2\rangle$ with probability 0, and $|j3\rangle$ with probability $\frac{2}{3}$. Of course, the sum of the probabilities adds up to 1.

Consider a system with

$$E = \{p_i; |j_i\rangle\}_{g_i} \quad (2)$$

with $\langle j_i | j_j \rangle = \delta_{ij}$. Suppose $i = 1$, then $E = \{p_1; |j_1\rangle; p_2; |j_2\rangle\}$. If we measure an observable \hat{A} in E , then

$$\begin{aligned} \langle \hat{A} \rangle_E &= \sum_j p_j \langle j | \hat{A} | j \rangle \\ &= \sum_j p_j \langle j | \hat{A} | j \rangle \\ &= p_1 \langle j_1 | \hat{A} | j_1 \rangle \end{aligned} \quad (3)$$

Now, what if $E = \{p_1; |j_1\rangle; p_2; |j_2\rangle\}$? If you think of p_1 and p_2 as giving weight to their corresponding state's expectation value, you can guess

$$\langle \hat{A} \rangle_E = p_1 \langle j_1 | \hat{A} | j_1 \rangle + p_2 \langle j_2 | \hat{A} | j_2 \rangle$$

and finally, if $E = \{p_i; |j_i\rangle\}_{g_i}$, then

$$\langle \hat{A} \rangle_E = \sum_i p_i \langle j_i | \hat{A} | j_i \rangle \quad (4)$$

Claim:

$$I = \sum_n |j\rangle\langle j| \tag{5}$$

where $\{|j\rangle\}_n$ is a complete set, i. e., $\sum_n |j\rangle\langle j| = I$, and I is the identity operator.

Proof:

$$I |j\rangle\langle j| = \sum_n |j\rangle\langle j| |j\rangle\langle j| = \sum_n |j\rangle\langle j| = |j\rangle\langle j|$$

QED

Thus,

$$\begin{aligned} \hat{A}^{D,E} &= \sum_i \rho_i |i\rangle\langle i| \hat{A} \\ &= \sum_i \rho_i |i\rangle\langle i| \sum_n |j\rangle\langle j| \hat{A} \\ &= \sum_i \rho_i \sum_n |i\rangle\langle j| \langle j| \hat{A} |i\rangle \\ &= \sum_i \rho_i \sum_n \langle j| \hat{A} |i\rangle |j\rangle\langle j| \\ &= \sum_i \rho_i \sum_n \langle j| \hat{A} |i\rangle |j\rangle\langle j| \\ &= \sum_n \left(\sum_i \rho_i \langle j| \hat{A} |i\rangle \right) |j\rangle\langle j| \end{aligned} \tag{6}$$

where

$$\hat{A} = \sum_i \rho_i |i\rangle\langle i| \tag{7}$$

is the density operator.

For some matrix B , we have $B_{ij} = \langle i|B|j\rangle$, thus $\hat{A}^{D,E} = \sum_n (\hat{A}^n)_{nn} = \text{Tr}(\hat{A})$, which gives the very useful and important result

$$\hat{A}^{D,E} = \text{Tr}(\hat{A}) \tag{8}$$

Note that the trace is independent of which complete set you use. Suppose $\{|j\rangle\}_n$ is a complete set, then using (5) we get

$$\begin{aligned} \text{Tr}(B) &= \sum_i \langle i|B|i\rangle = \sum_i \langle i|B|j\rangle \langle j|i\rangle \\ &= \sum_i \langle i|B|j\rangle \sum_n |j\rangle\langle j| \\ &= \sum_j \langle j|i\rangle \langle i|B|j\rangle \\ &= \sum_j \langle j|B|j\rangle \end{aligned} \tag{9}$$

thus, just think of the complete basis $\{|j\rangle\}_n$, where for i , the i^{th} component is 1 and the other components are 0, to build your intuition.

Claim 2: $Tr(\hat{\rho}) = 1$

Proof:

$$\begin{aligned}
 Tr(\hat{\rho}) &= \sum_n \langle n | \hat{\rho} | n \rangle = \sum_n \langle n | \sum_i p_i | i \rangle \langle i | n \rangle \\
 &= \sum_i p_i \sum_n \langle n | i \rangle \langle i | n \rangle \\
 &= \sum_i p_i \langle i | i \rangle \\
 &= \sum_i p_i = 1
 \end{aligned} \tag{20}$$

QED

where we used the normalization property of the wavefunction, and (5).

Claim 3: $\hat{\rho}$ can be diagonalized to the following

$$\hat{\rho} = \begin{pmatrix} p_1 & 0 & \dots & 0 \\ 0 & p_2 & 0 & \dots & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & \dots & \dots & \dots & p_j \end{pmatrix} \tag{21}$$

where j is the size of the set.

Proof:

We have that $\hat{\rho} = \sum_i p_i | i \rangle \langle i |$, thus

$$\langle j | \hat{\rho} | k \rangle = \sum_i p_i \langle j | i \rangle \langle i | k \rangle = p_j \delta_{jk} \tag{22}$$

so $| i \rangle$ as eigenvectors $f_j | i \rangle g_k$ with eigenvalues $f_j p_k g_k$. Let

$$M = \begin{pmatrix} | j_1 \rangle & | j_2 \rangle & \dots & | j_j \rangle \end{pmatrix} \tag{23}$$

Then

$$\hat{\rho} M = \begin{pmatrix} | j_1 \rangle & | j_2 \rangle & \dots & | j_j \rangle \end{pmatrix} \begin{pmatrix} p_1 & 0 & \dots & 0 \\ 0 & p_2 & 0 & \dots & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & \dots & \dots & \dots & p_j \end{pmatrix} = M \begin{pmatrix} p_1 & 0 & \dots & 0 \\ 0 & p_2 & 0 & \dots & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & \dots & \dots & \dots & p_j \end{pmatrix} \tag{24}$$

which we can rewrite in a more enlightening way

$$\hat{\rho} M = \begin{pmatrix} | j_1 \rangle & | j_2 \rangle & \dots & | j_j \rangle \end{pmatrix} \begin{pmatrix} p_1 & 0 & \dots & 0 \\ 0 & p_2 & 0 & \dots & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & \dots & \dots & \dots & p_j \end{pmatrix} = M \begin{pmatrix} p_1 & 0 & \dots & 0 \\ 0 & p_2 & 0 & \dots & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & \dots & \dots & \dots & p_j \end{pmatrix} \tag{25}$$

if M is invertible, we get

3 The Canonical Ensemble

3.1 Lagrange Multipliers

Let S be a surface given by $f(x; y; z) = c$ for some constant c . What if you want to find which point on this surface is the closest to the origin? To do this we need some background. A curve on this surface is defined by $r(t) = (x(t); y(t); z(t))$. Now, the derivative of this curve, $r'(t)$, is tangent to the curve at any point $P = (x(t); y(t); z(t))$. Let's take one of them $P_0 = (x(0); y(0); z(0)) = (x_0; y_0; z_0)$. At P_0 , $r'(t)$ is tangent to the curve and thus tangent to S , which means it lies on the tangent plane of $f(x; y; z)$ at P_0 . This means that if a vector is perpendicular to $r'(t)$ at P_0 , then it is perpendicular to the tangent plane of S at P_0 .

Claim: The gradient of f is perpendicular to $r'(t)$ at any point.

If we can prove that $r' \cdot \nabla f|_{P_0} = 0$ for some arbitrary point P_0 we would be done. But

$$\frac{df}{dt}|_{P_0} = \frac{\partial f}{\partial x}|_{P_0} \frac{dx}{dt}|_{t_0} + \frac{\partial f}{\partial y}|_{P_0} \frac{dy}{dt}|_{t_0} + \frac{\partial f}{\partial z}|_{P_0} \frac{dz}{dt}|_{t_0} = r' \cdot \nabla f|_{P_0} \tag{30}$$

but $f(x; y; z) = c$ where c is constant, thus $\frac{df}{dt}|_{P_0} = 0$.

QED

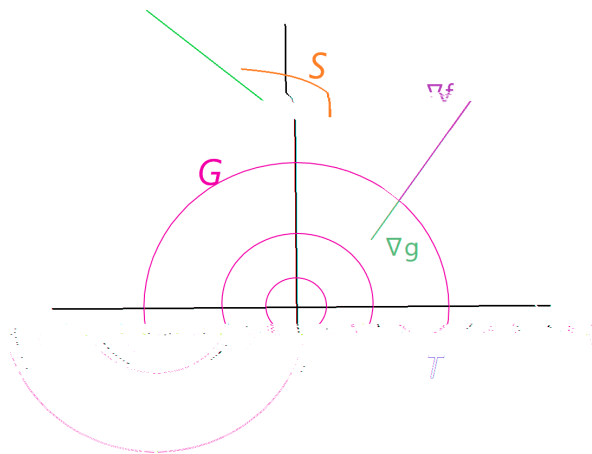


FIG. 1. Lagrange Multipliers

Let's go back to our minimization problem. Try the problem in two dimension, where S is a curve. Take a compass and start drawing circles with a common center at the origin. Start from a small one and increase the radius of the next one. Continue increasing the radius until a circle enters in contact with S

where g is the function containing the constraint, and λ is the famous Lagrange multipliers. The general case is given by

$$r f = \sum_i \lambda_i g_i \quad (32)$$

3.2 Mixed States and Entropy

We know that the entropy of a system is given by

$$S = -k \sum_i p_i \ln p_i = -k \text{Tr}(\rho \ln \rho) \quad (33)$$

ρ is diagonalized. What are the $f p_i g$ in the minimum entropy state? From (33), it's easy to see that S is a of positive quantities $\Rightarrow S \geq 0 \Rightarrow S_{min} = 0$. Recall that we need $\sum_i p_i = 1$, which means that $p_k = 1$ and $p_i = 0 \forall i \neq k$. This might ring a bell (no pun intended), we have in this case a pure state. Thus, in order to have minimum entropy, an ensemble must be composed of a pure state.

Now what about maximizing entropy? This is where Lagrange multipliers come handy. In this case, $f = S$, and $g = \sum_j p_j = 1$. Looking back at (32), we must have

$$-k \sum_i p_i \ln p_i = \lambda \sum_i p_i = \lambda \quad (34)$$

We can simplify this quite a bit

$$\begin{aligned} -k \sum_i p_i \ln p_i &= \lambda \sum_i p_i \\ &= \sum_i (k \ln p_i + \lambda) p_i \end{aligned} \quad (35)$$

but λ is arbitrary, thus

$$k \ln p_i + \lambda = 0 \Rightarrow p_i = e^{-\lambda/k} \quad (36)$$

Using our constraint $\sum_i p_i = 1$, we get

$$\sum_{i=1}^j e^{-\lambda/k} = 1 \Rightarrow e^{-\lambda/k} = \frac{1}{j} \Rightarrow p_i = \frac{1}{j} \quad (37)$$

This means that in order to maximize entropy, we want a uniform distribution. This gives us that a state is a maximally mixed state if every eigenvalue of the density operator is equal, i.e.,

$$\rho = \frac{1}{j} \begin{pmatrix} 1 & 0 & \dots & 0 \\ 0 & 1 & 0 & \dots & 0 \\ \vdots & & & & \vdots \\ 0 & \dots & 0 & \dots & 1 \end{pmatrix}$$

This gives us

3.3 Statistical Mechanics

To some extent, statistical mechanics is an assumption about the density matrix for a macroscopic system. The assumptions (constraints) are

- (a) $\sum_i p_i = 1$.
- (b) $\hat{H} = E$ is known.

Constraint (b) can be expressed in a more useful way since from Schroedinger equation $\hat{H} |j\rangle = E_j |j\rangle$ and $\langle j | \hat{H} |k\rangle = E_k \delta_{jk}$, which gives

$$\hat{H} = \sum_k p_k \hat{H} |k\rangle \langle k| = \sum_k p_k E_k = E \quad (40)$$

Since we have two constraints, using equation (32) we have that

$$\sum_i p_i \ln p_i = - \sum_i p_i \ln p_i + \sum_i p_i E_i = -E \quad (41)$$

which we can simplify to

$$\sum_i p_i \ln p_i + \sum_i p_i \ln p_i = - \sum_i p_i E_i = -E \quad (42)$$

but p_i is arbitrary, thus

$$\ln p_i + E_i = 0 \Rightarrow p_i = e^{-\frac{E_i}{k}} \quad (43)$$

Using our constraint $\sum_i p_i = 1$, we get

$$e^{-\frac{E_i}{k}} = \frac{1}{Z} \quad (44)$$

with

$$Z = \sum_i e^{-\frac{E_i}{k}} \quad (45)$$

Letting $\frac{1}{k} = \beta$, we get the well known canonical ensemble equations

$$p_i = \frac{e^{-\beta E_i}}{Z} \quad (46)$$

and

$$Z = \sum_i e^{-\beta E_i} \quad (47)$$

4 Quantum Mechanics and Statistical Mechanics

4.1 Mixed State, Pure State, and Temperature

The internal energy of a monoatomic gas is given by $E = \frac{3}{2}NkT \propto E \propto T$, but looking back at (47), we must have $\propto \frac{1}{E}$ since the exponential must be unitless. These two relations imply that $\propto \frac{1}{T}$. Thus,

$$T \rightarrow 0^+ \Rightarrow Z = \sum_j p_j \Rightarrow p_i = \frac{1}{j}$$

This means that all states are equally probable, which is what we found for the maximally mixed state. What about $T \rightarrow 0^+$? I'm sure you can guess what is about to happen

$$Z = \sum_{i=0}^{\infty} e^{-E_i} = e^{-E_0} \sum_{i=0}^{\infty} e^{-(E_i - E_0)} = e^{-E_0} \sum_{i=1}^{\infty} e^{-(E_i - E_0)} \quad (48)$$

therefore since $E_i > E_0 \forall i > 0$, we get

$$T \rightarrow 0^+ \Rightarrow \sum_{i=1}^{\infty} e^{-(E_i - E_0)} \rightarrow 0 \Rightarrow Z \rightarrow e^{-E_0} \quad (49)$$

this means that

$$T \rightarrow 0^+ \Rightarrow \sum_{i=1}^{\infty} p_k = \frac{e^{-E_k}}{Z} \rightarrow 0 \quad \forall k \neq 0 \quad (50)$$

thus, $p_k = 0 \forall k \neq 0$ and $p_0 = 1$. Therefore all the particles are in the ground state E_0 . This is exactly what we found for a pure state.

4.2 Imaginary Statistical Mechanics

Definition: (Function of a Hermitian Operator) Suppose that a Hermitian operator A has a spectral decomposition $A = \sum_i a_i |i\rangle\langle i|$ for some orthonormal basis $|i\rangle$. Then the operator $f(A)$ for some function f is defined as follows:

In quantum mechanics, we can determine how a quantum state evolves with time using the operator $U(t)$.

$$|j(t)\rangle = U(t)|j(0)\rangle \quad (56)$$

We can get it from looking at Schrodinger equation.

$$\hat{H}|j(t)\rangle = i\hbar \frac{\partial}{\partial t}|j(t)\rangle \quad |j(t)\rangle = e^{-\frac{i\hat{H}t}{\hbar}}|j(0)\rangle \quad (57)$$

thus, $U(t) = e^{-\frac{i\hat{H}t}{\hbar}}$.

By making the change

$$t \rightarrow i\hbar \quad (58)$$

we get

$$U(t) = U(i\hbar) = e^{-\hat{H}} \quad (59)$$

which gives us the startling result that

whic

$$\text{Tr}[U(i\hbar)] = \text{Tr} e^{-\hat{H}}$$

but then $T^0 = \cos(z \tan + T)$. We can simplify this by using the fact that $\sec^2 + \tan^2 = 1$
we get $\sec = \frac{1}{\sqrt{1 - \tan^2}} = \frac{1}{\cos}$. Thus